

# K-Stage Bipartite Matching

- Inspired by applications in matching market



- K-stage bipartite matching (with batch arrival)
  - offline supply nodes
  - K stages, new batch of demand nodes arrives at each stage
  - make irrevocable matching decision
- goal: maximize the competitive ratio against optimum in-hindsight solution

$$\Gamma(K) \triangleq \sup_{\text{ALG}} \left( \inf_{G, \{w_j\}_{j \in V}} \frac{\mathbb{E} [\text{weight}(\text{ALG})]}{\text{weight}(\text{OPT-OFFLINE})} \right)$$

# K-Stage Bipartite Matching



$K = 1$   
(fully offline)

$$\Gamma(1) = 1$$



$K > 1$   
(multi-stage)



$K = |U|, \text{batch-size}=1$   
(fully online)

$$\Gamma(\infty) = 1 - \frac{1}{e}$$

## [Formal Research Question]

In the middle-ground regime with  $K > 1$  batches, what is the optimal competitive ratio achieved by a multi-stage algorithm?

$$\text{Main Result: } \Gamma(K) = 1 - \left(1 - \frac{1}{K}\right)^K$$

### Theorem

**[Lower-bound]** For any  $K > 0$ , we propose an efficient (*simple*) multi-stage algorithm with competitive ratio no smaller than  $\Gamma(K) = 1 - \left(1 - \frac{1}{K}\right)^K$ .

**[Upper-bound]** For any  $K > 0$ , there exists an instance for which no multi-stage algorithm has a competitive ratio better than  $\Gamma(K) = 1 - \left(1 - \frac{1}{K}\right)^K$ .

We show same results hold when:

- *Fractional*: Vertex-weighted b-matching, budgeted allocation, assortment
- *Integral*: above settings, with the large budget assumption

# Attacking the Problem from a New Perspective

- **“Polynomial regularized convex-programming matching”**
  - Identifying a particular **sequence of polynomials of decreasing degrees**, one for each batch
  - Solving resulting **regularized max-weight matching convex programs** to find a family of balanced allocations
  - Using **graph decompositions** induced by optimal solutions to come up with a new **primal-dual competitive analysis**

